

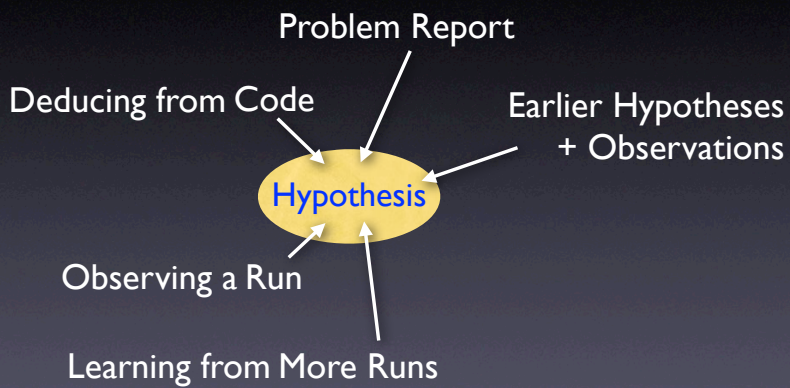
# Deducing Errors

Andreas Zeller



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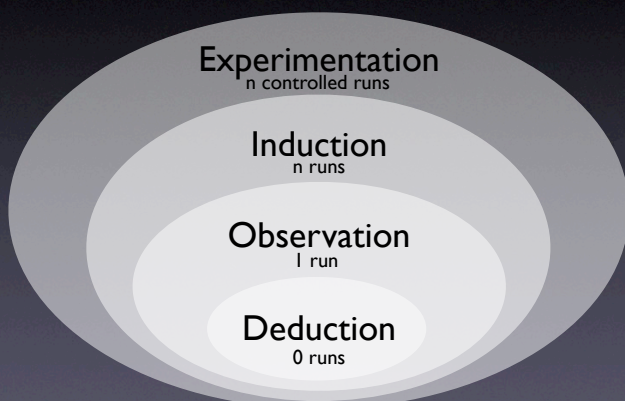
# Obtaining a Hypothesis



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# Reasoning about Runs



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# Reasoning about Runs

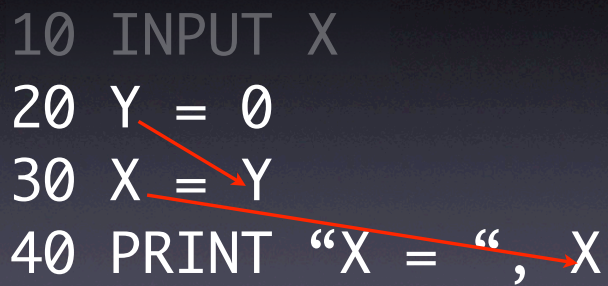
Deduction  
0 runs

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## What's relevant?

```
10 INPUT X
20 Y = 0
30 X = Y
40 PRINT "X = ", X
```



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## Fibonacci Numbers

$$fib(n) = \begin{cases} 1, & \text{for } n = 0 \vee n = 1 \\ fib(n-1) + fib(n-2), & \text{otherwise} \end{cases}$$

1	1	2	3	5	8	13	21	34	55
---	---	---	---	---	---	----	----	----	----

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# fibonacci.c

```
int fib(int n)
{
    int f, f0 = 1, f1 = 1;

    while (n > 1) {
        n = n - 1;
        f = f0 + f1;
        f0 = f1;
        f1 = f;
    }

    return f;
}

int main()
{
    int n = 9;


    while (n > 0)
    {
        printf("fib(%d)=%d\n",
              n, fib(n));
        n = n - 1;
    }

    return 0;
}
```

## Fibo in Action

```
$ gcc -o fibo fibonacci.c
$ ./fibo
fib(9)=55
fib(8)=34
...
fib(2)=2
fib(1)=134513905
```

Where does fib(1) come from?



## Effects of Statements

- **Write.** A statement can change the program state (i.e. write to a variable)
- **Control.** A statement may determine which statement is executed next (other than unconditional transfer)

# Affected Statements

- **Read.** A statement can read the program state (i.e. from a variable)
- **Execution.** To have any effect, a statement must be executed.

# Effects in fibo.c

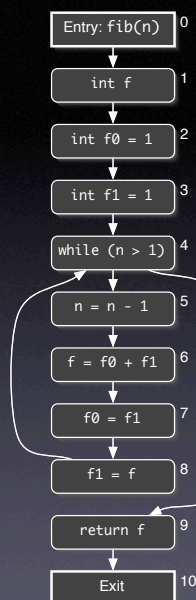
	Statement	Reads	Writes	Controls
0	fib(n)		n	1-10
1	int f		f	
2	f0 = 1		f0	
3	f1 = 1		f1	
4	while (n > 1)	n		5-8
5	n = n - 1	n	n	
6	f = f0 + f1	f0, f1	f	
7	f0 = f1	f1	f0	
8	f1 = f	f	f1	
9	return f	f	<ret>	

# Control Flow

```
int fib(int n)
{
    int f, f0 = 1, f1 = 1;

    while (n > 1) {
        n = n - 1;
        f = f0 + f1;
        f0 = f1;
        f1 = f;
    }

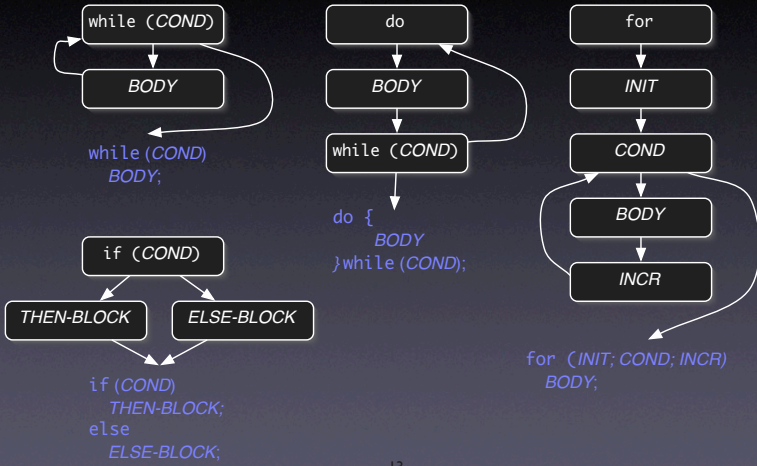
    return f;
}
```



The CFG is best developed incrementally on an extra board.



# Control Flow Patterns

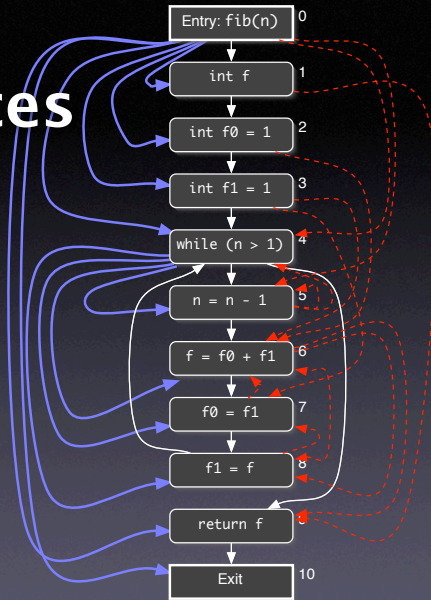


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# Dependences

- A - - - -> B**  
**Data dependency:**  
 A's data is used in B;  
 B is data dependent on A
- A . . . .> B**  
**Control dependency:**  
 A controls B's execution;  
 B is control dependent on A



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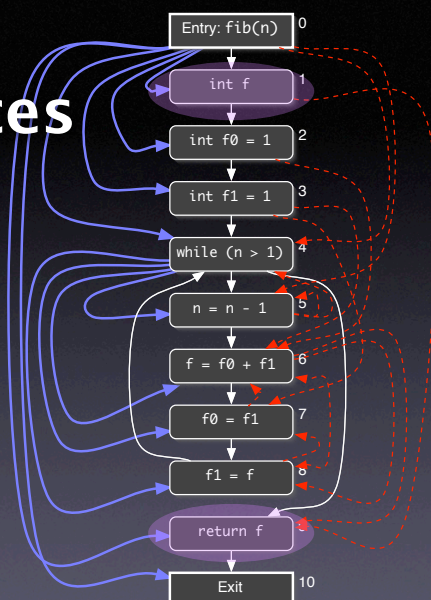
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Again, this is best developed interactively on the board (possibly by having the students call further dependences)

# Dependences

Following the dependences, we can answer questions like

- Where does this value go to?
- Where does this value come from?

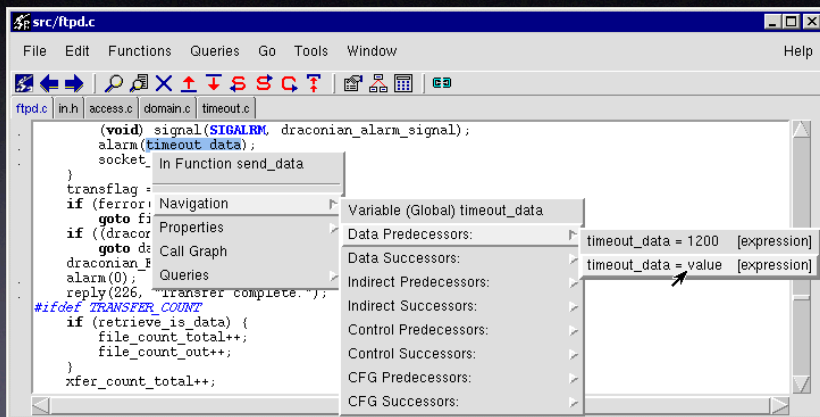


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Again, this is best developed interactively on the board (possibly by having the students call further dependences)

# Navigating along Dependences



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# Program Slicing

- A *slice* is a subset of the program
- Allows programmers to *focus on what's relevant* with respect to some statement *S*:
  - All statements influenced by *S*
  - All statements that influence *S*

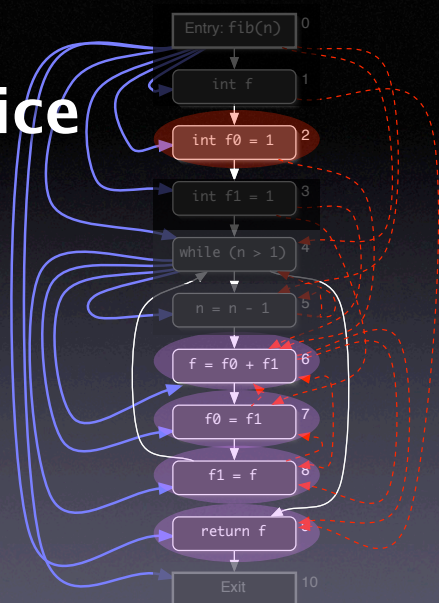
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# Forward Slice

- Given a statement *A*, the forward slice contains all statements whose read variables or execution could be influenced by *A*
- Formally:  

$$S^F(A) = \{B | A \rightarrow^* B\}$$



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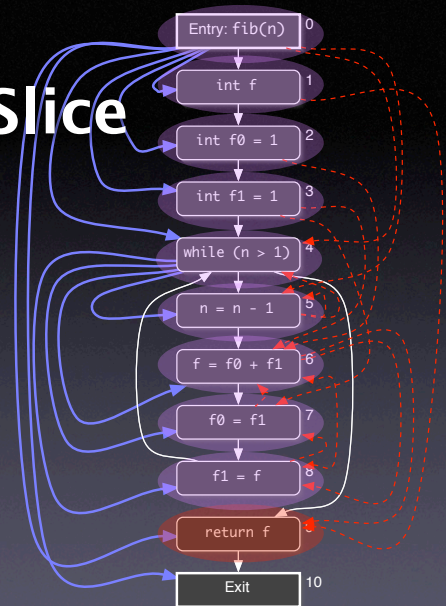
Again, this is best developed interactively on the board (possibly by having the students call further dependences)



# Backward Slice

- Given a statement B, the backward slice contains all statements that could influence the read variables or execution of B
- Formally:  

$$S^B(B) = \{A \mid A \rightarrow^* B\}$$



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Again, this is best developed interactively on the board (possibly by having the students call further dependences)

# Two Slices

```
int main() {
  int a, b, sum, mul;
  sum = 0;
  mul = 1;
  a = read();
  b = read();
  while (a <= b) {
    sum = sum + a;
    mul = mul * a;
    a = a + 1;
  }
  write(sum);
  write(mul);
}
```

Slice Operations:

- Backbones
- Dices
- Chops

← Backward slice of sum  
 ← Backward slice of mul

20

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# Backbone

```
a = read();
b = read();
while (a <= b) {
  a = a + 1;
}
```

- Contains only those statements that occur in both slices
- Useful for focusing on common behavior

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# Two Slices

```
int main() {  
    int a, b, sum, mul;  
    sum = 0;  
    mul = 1;  
    a = read();  
    b = read();  
    while (a <= b) {  
        sum = sum + a;  
        mul = mul * a;  
        a = a + 1;  
    }  
    write(sum);  
    write(mul);  
}
```

Slice Operations:

- Backbones
- Dices
- Chops

← Backward slice of sum

← Backward slice of mul

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# Dice

```
sum = 0;
```

```
sum = sum + a;
```

```
write(sum);
```

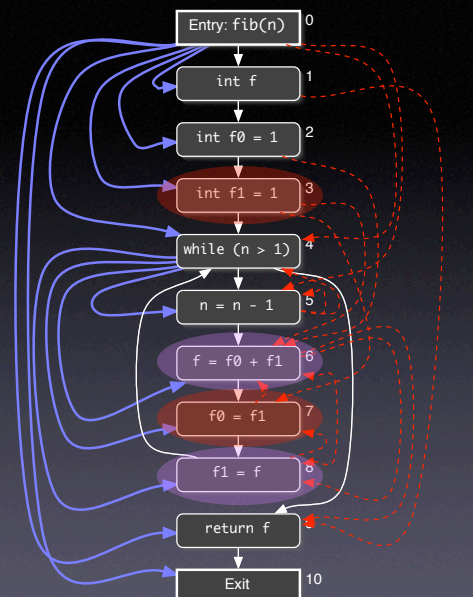
- Contains only the difference between two slices
- Useful for focusing on differing behavior

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# Chop

- Intersection between a forward and a backward slice
- Useful for determining influence paths within the program



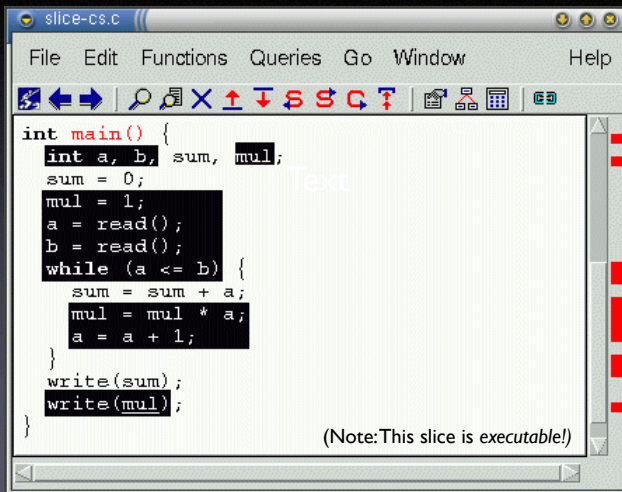
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Again, this is best developed interactively on the board (possibly by having the students call further dependences)



# Leveraging Slices



```
int main() {
    int a, b, sum, mul;
    sum = 0;
    mul = 1;
    a = read();
    b = read();
    while (a <= b) {
        sum = sum + a;
        mul = mul * a;
        a = a + 1;
    }
    write(sum);
    write(mul);
}
```

(Note: This slice is executable!)

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## Deducing Code Smells

- Use of uninitialized variables
- Unused values
- Unreachable code
- Memory leaks
- Interface misuse
- Null pointers

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## Uninitialized Variables

```
$ gcc -Wall -O -o fibo fibo.c
fibo.c: In function `fib':
fibo.c:7: warning: `f' might be
used uninitialized in this
function
```

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# False Positives

```
int go;
switch (color) {
  case RED:
  case AMBER:
    go = 0;
    break;
  case GREEN:
    go = 1;
    break;
}
if (go) { ... }
```

warning: `go` might be used uninitialized in this function

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# Unreachable Code

```
if (w >= 0)
  printf("w is non-negative\n");
else if (w > 0)
  printf("w is positive\n");
```

warning: will never be executed

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# Memory Leaks

```
int *readbuf(int size)
{
  int *p = malloc(size * sizeof(int));
  for (int i = 0; i < size; i++) {
    p[i] = readint();
    if (p[i] == 0)
      return 0; // end-of-file
  }
  return p;
}
```

memory leak

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# Interface Misuse

```
void readfile()
{
    int fp = open(file);
    int size = readint(file);
    if (size <= 0)
        return;
    ...
    close(fp);
}
```

stream not closed

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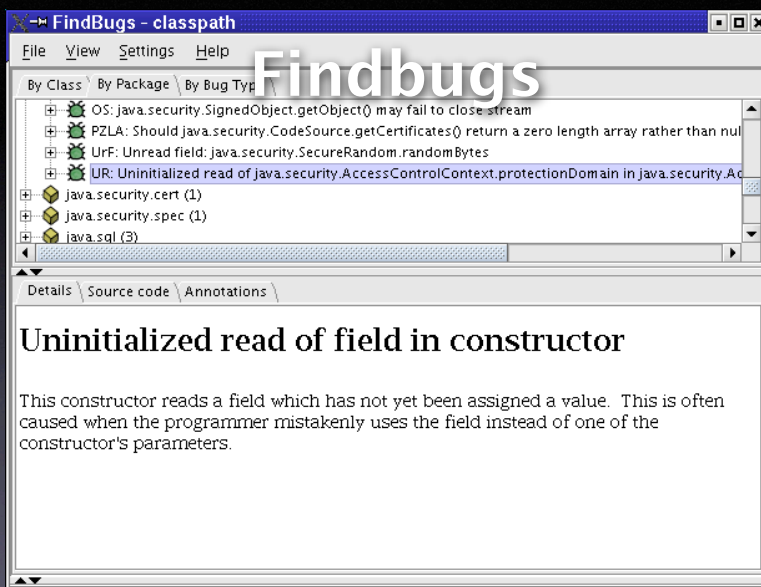
# Null Pointers

```
int *readbuf(int size)
{
    int *p = malloc(size * sizeof(int));
    for (int i = 0; i < size; i++) {
        p[i] = readint();
        if (p[i] == 0)
            return 0; // end-of-file
    }
    return p;
}
```

p may be null

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The screenshot shows the FindBugs application window titled "FindBugs - classpath". The main window is divided into two panes. The top pane, titled "By Class", lists several classes with bug counts: "java.security.cert (1)", "java.security.spec (1)", and "java.sql (3)". The bottom pane, titled "Details", shows the details for a specific bug: "UR: Uninitialized read of field in constructor". The bug description reads: "This constructor reads a field which has not yet been assigned a value. This is often caused when the programmer mistakenly uses the field instead of one of the constructor's parameters."

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- Class implements Cloneable but does not define or use clone method
- Method might ignore exception
- Null pointer dereference in method
- Class defines equal(); should it be equals()?
- Method may fail to close database resource
- Method may fail to close stream
- Method ignores return value
- Unread field
- Unused field

## Defect Patterns

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## Limits of Analysis

```
int x;
for(i=j=k=1;--j||k;k=j?i%j?k:k-j:(j=i+=2));
write(x);
```

- Is x being used uninitialized or not?
- Loop halts only if there is an odd perfect number (= a number that's the sum of its proper positive divisors)
- Problem is undecidable yet

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```
static void shell_sort(int a[], int size)
{
    int i, j;
    int h = 1;
    do {
        h = h * 3 + 1;
    } while (h <= size);
    do {
        h /= 3;
        for (i = h; i < size; i++)
        {
            int v = a[i];
            for (j = i; j >= h && a[j - h] > v; j -= h)
                a[j] = a[j - h];
            if (i != j)
                a[j] = v;
        }
    } while (h != 1);
}
```

Conservative approximation:  
any a[] depends on all a[]

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# Causes of Imprecision

- Indirect access, as in `a[i]`
- Pointers
- Functions
- Dynamic dispatch
- Concurrency

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# Risks of Deduction

- **Code mismatch.** Is the run created from this very source code?
- **Imprecision.** A slice typically encompasses 90% of the source code.
- **Abstracting away.** Failures may be caused by a defect in the environment.

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# Dijkstra's Curse

Testing can only find the  
*presence of errors,*  
not their *absence*

configurations →

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But still, testing suffers from what I call Dijkstra's curse – a double meaning, as it applies both to testing as to his famous quote. Is there something that can find the **absence** of errors?

# Formal Verification



configurations

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# Formal Verification

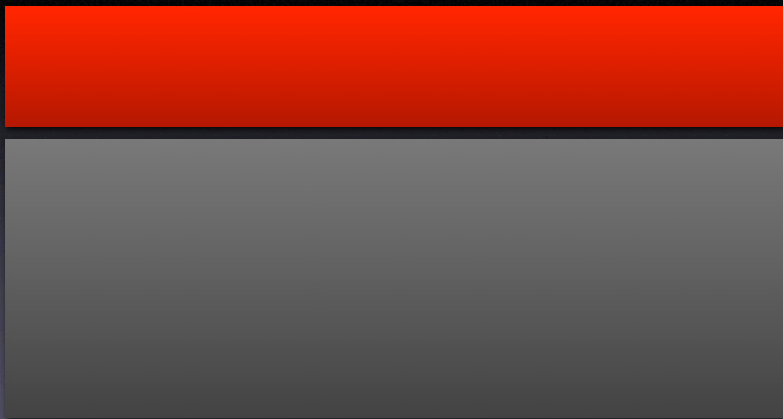


configurations

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# Formal Verification



configurations

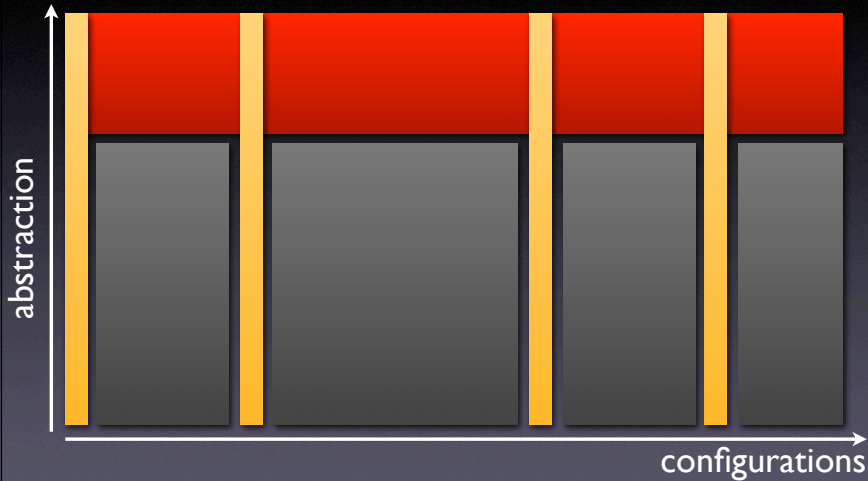
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Areas missing might be:  
the operating system, the  
hardware, all of the world  
the system is embedded in  
(including humans!)



# Best of Both Worlds



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We might not be able to cover **all** abstraction levels in **all** configurations, but we can do our best to cover as much as possible.

# Hetzel-Myers Law

A combination of different V&V methods outperforms any single method alone.

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# Increasing Precision

- **Verification.** If we know that certain properties hold, we can leverage them in our inference process.
- **Observation.** Facts from concrete runs can be combined with deduction.

...in the weeks to come!

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# Concepts

- ★ To reason about programs, use
  - deduction (0 runs)
  - observation (1 run)
  - induction (multiple runs)
  - experimentation (controlled runs)

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## Concepts (2)

- ★ To isolate value origins, follow back the dependences
- ★ Dependences can uncover *code smells* such as
  - uninitialized variables
  - unused values
  - unreachable code
- ★ Get rid of smells before debugging

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## Concepts (3)

- ★ To slice a program, follow dependences from a statement *S* to find all statements that
  - could be influenced by *S* (forward slice)
  - could influence *S* (backward slice)

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# Concepts (4)

- ★ Using deduction alone includes a number of risks, including code mismatch, sbstracting away, and imprecision.
- ★ Any deduction is limited by the halting problem and must thus resort to conservative approximation.
- ★ For debugging, deduction is best combined with actual observation.

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